

BGP Security Vulnerabilities Analysis

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Agenda

- 1. Introduction
- 2. Possible Attacks
- 3. Vulnerabilities and Risks
- 4. Security Considerations
- 5. References



Introduction



Introduction: History

- BGP 4 (RFC 1771) specified in March 95
- BGP 3 (RFC 1267) specified in October 91
- Based on EGP (RFC 904) of April 84
- BGP was created when the Internet was much more peaceful than nowadays
- It lacks protection against errors and authentication

Introduction

BGP is a TCP/IP – protocol

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- Subject to TCP/IP attacks like IP Spoofing, Session stealing, etc.
- Outsiders could inject bogus routing information or disrupt peer to peer communication
- This new information would spread through peers
- Therefore at least authentication mechanism must be supported (TCP MD5 Signature)

Introduction

 Faulty routing information can be caused by misconfigured peers themselves

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- By masquerading as other legitimate BGP speakers
- By distributing unauthorized routing information
- Whole portions of the network could become unreachable
- Packets could be forwarded by a suboptimal path or a path that will not forward the traffic
- Therefore traffic could be delayed or misleaded

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Introduction

- The damage resulting from attacks might be:
 - Starvation
 - Network congestion
 - Blackhole
 - Delay
 - Looping
 - Eavesdrop
 - Partition
 - Cut
 - Churn
 - Instability
 - Overload
 - Resource exhaustion



BGP is subject to the following attacks:

- Eavesdropping:
 - Routing data is carried in cleartext (attacks confidentiality)
- Replay:

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- BGP doesn't provide any protection against replay attacks
- Message Insertion:
 - No protection against message insertion
 - However if TCP Session is fully established, prediction of the correct session number becomes necessary for the attacker

- Further attacks might be:
 - Message deletion:

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- No protection inside BGP
- Again difficult against mature TCP implementation
- Message modification:
 - Modifications not altering the length of the payload can not be detected
- Man-in-the-middle:
 - As BGP has no peer entity authentication, man-in-the-middle attacks are easy to accomplish

- Another attack against BGP is the Denial of service attack:
 - Bogus routing data can represent a DoS attack to:
 - End systems trying to transmit data through the network
 - The network infrastructure itself

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- Certain bogus information can represent a DoS attack to the BGP protocol itself:
 - E.g.: advertising large numbers of more specific routes can cause BGP traffic and routing table size to explode

Attacks: Countermeasures

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- The protection of BGP using the TCP MD5 signature option (RFC 2385) will counter most of the previously listed attacks from outsiders
- It will not protect against eavesdropping, but confidentiality of routing data is no goal of BGP
- Replay attacks will still be possible too, but with TCP sequence number processing it will be hard to accomplish
- Still no protection against misconfigured legitimate speakers





- There are three major vulnerabilities in BGP:
 - There is no mechanism to proof freshness, protection of integrity and peer authentication in the BGP protocol
 - There is no validation of the authority of an Autonomous System (AS) to announce Network Layer Reachability Information (NLRI)
 - There is no insurance of the authenticity of path attributes announced by an AS



- The first of these vulnerabilities motivated the support of the TCP MD5 signature in the BGP specification
- If implemented correctly, it provides message integrity and peer authentication
- But in the spec. the MD5 algorithm is supposed to be secure (which is not true), and that the shared secret is protected and difficult to guess



- There are four diffrent types of BGP messages:
 - OPEN
 - KEEP ALIVE
 - NOTIFICATION
 - UPDATE
- Each of them has ist own vulnerabilities, which will be, besides other vulnerabilities, discussed in the following



- Message Header:
 - Each BGP message starts with a standard header
 - Sytactic errors within the header will cause the connection to be closed, newly learned routes will be deleted and a new decision process about routes will be started



- OPEN message:
 - Receipt of an OPEN message in state Connect, Active or Established, or receipt of erroneous OPEN messages will cause:
 - Closing of connection
 - Deletion of all associated routes
 - Starting of decision process
 - Return state to idle
 - Receipt of an OPEN message in state OpenSent (spoofing) will cause transition to OpenConfirmed state and the following legitimate OPEN message will be dropped



- KEEPALIVE message:
 - Receipt of a KEEPALIVE message when the peering connection is in the Connect, Active or OpenSent state would cause a transition to the Idle state, and the failing of the connection to be established
 - To exploit this vulnerability, the KEEPALIVE message must be timed carefully within the exchanged messages



- Receipt of a NOTIFICATION message in any state will cause the previosly described effects:
 - Closing of connection
 - Deletion of all associated routes
 - Starting of decision process
 - Return state to idle



- UPDATE message:
 - In general, the UPDATE message carries the routing information, therefore the ability to spoof any part of this message will alter the routing tables
 - Withdrawn Routes field inside an UPDATE message:
 - By modifying this field an attacker could cause the elimination of existing legitimate routes
 - Reestablished routes could be deleted via replaying a previously recorded withdrawal
 - But the withdrawal of routes can only be performed by the BGP speaker having formerly announced these routes



- UPDATE message continued:
 - The Path Attributes within the UPDATE message present various vulnerabilities and risks:
 - Altering of the AS_PATH attribute could be used to affect routing decisions, and thus mislead traffic to suboptimal routes, to create loops or to gain access to traffic
 - The NEXT_HOP attribute could be modified to disrupt forwarding of traffic between to AS's, or to force another AS to carry traffic it would otherwise not have to



- UPDATE message continued:
 - Modifying or forging the NLRI field in the UPDATE message could cause :
 - Disruption of routing to the announced network
 - Overwhelming of a router along the announced route
 - Data loss if the announced route will not forward traffic to the announced network
 - Routing of traffic by a suboptimal route, etc.
 - In general, syntactic malformed UPDATE messages will cause the connection to be closed, associated routes will be deleted, etc., with the previosly described effects



- Other vulnerabilities arise through the use of the TCP protocol:
 - TCP SYN attack:
 - BGP is vulnerable to SYN flooding as other protocols using TCP
 - An attacker could send a SYN, and a sequence of BGP packets to establish a BGP session, letting the legitimate connection appear as a collision which would be destroyed
 - TCP SYN ACK:
 - If an attacker could answer to a SYN before the legitimate peer, which would receive an empty ACK reply this would finally result in a RST that would break the connection



- Further spoofed RST or FIN messages would also cause the connection to be broken
- All these TCP attacks can be countered by the use of BGP session protection via the TCP MD5 signature option
- DoS and DDoS attacks against BGP are easy to accomplish, because packets directed to port 179 are passed to the BGP process, normally residing on a slower processor



Security Considerations



Security Considerations

- The use of the ,Protection of BGP Sessions via the TCP MD5 Signature Option' (RFC2385) counters message insertion, message deletion, modification and man-in-the-middle attacks from outsiders and therefore should be used
- If routing data confidentiality is desired, this could be accomplished using IPSec ESP
- Both provide security, assuming the algorithms are secure, the used secrets are protected from exposure and not guessable, the platforms are secure, etc.



Security Considerations Residual Risks

- Protection against attacks arising from legitimate peers could be accomplished through:
 - Origination Protection: sign the originating AS
 - Origination and Adjacency Protection: sign the originating AS and predecessor information
 - Origination and Route Protection: sign the originating AS and remove AS_PATHs of ,bad routers' (Secure-BGP)
 - Filtering: verify AS_PATH and NLRI originating AS via a registry (RFC2725)
- Except of Filtering, which is limited to the ,outscirts' of the internet, none of these is in common use

Security Considerations Operational Protections

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- BGP is used by all major ISPs, to distribute global routing information, internally and between each other
- Therefore BGP implementations are confronted with huge amounts of traffic, making use of cryptography nearly impossible
- Protection against DoS attacks can only be achieved using port based packet filtering

Security Considerations Operational Protections

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- Current practice of the ISPs is the usage of filtering techniques at their borders, reducing exposure to attacks from outside
- These filters remove the BGP Port Number (179) from traffic destined to the inside, preventing internal peers to be flooded
- Prevented from injecting sufficent traffic from the outside, attackers have to gain physical access

References

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Thank You!